# Computer Algebra for Lattice Path Combinatorics

Alin Bostan

based on joint works with F. Chyzak, M. Van Hoeij, M. Kauers, L. Pech, K. Raschel, B. Salvy

> Séminaire de Combinatoire Philippe Flajolet, Institut Henri Poincaré, March 28, 2013

## Why Lattice Paths?

Applications in many areas of science

- probability theory (branching processes, games of chance, ...)
- operations research (queueing theory, ...)
- discrete mathematics (permutations, trees, words, urns, ...)
- statistical physics (lsing model, ...)

Journal of Statistical Planning and Inference 140 (2010) 2237-2254



#### A history and a survey of lattice path enumeration

#### Katherine Humphreys

Department of Mathematical Sciences, Florida Atlantic University, Boca Raton, FL 33431, USA

| ARTICLE INFO                                      | A B S T R A C T   |
|---|---|
| Available online 21 January 2010                  | In celebration of the Sixth International Conference on Lattice Path Counting and   |
| Keywords:<br>Lattice path<br>Reflection principle | Applications, it is befitting to review the history of lattice path enumeration and to<br>survey how the topic has progressed thus far.<br>We start the history with early games of chance specifically the ruin problem which  |
| Method of images                                  | later appears as the ballot problem. We discuss André's Reflection Principle and its<br>misnomer, its relation with the method of images and possible origins from physics and<br>Brownian motion, and the earliest evidence of lattice path techniques and solutions.<br>In the survey, we give representative articles on lattice path techniques and solutions.<br>The literature in the last 35 years by the lattice, step set, boundary, characteristics |
|   | counted, and solution method. Some of this work appears in the author's 2005<br>dissertation.   |

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# Why Computer Algebra?

Because we like it!







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Discrete Mathematics 306 (2006) 992-1021

#### DISCRETE MATHEMATICS

THE EVOLUTION OF TWO STACKS IN BOUNDED SPACE AND RANDOM WALKS IN A TRIANGLE

#### Philippe FLAJOLET

INRIA Rocquencourt 78150 Le Chesnay (France)

#### ABSTRACT

#### Combinatorial aspects of continued fractions

P. Flajolet

IRIA, 78150 Rocquencourt, France

#### Abstract

We show that the universal continued fraction of the Stielije-Jacobi yep is equivalent to the characteristic series of labelled paths in the plane. The equivalence holds in the soft series in non-commutative indeterminants. Using it, we derive direct combinatorial proofs of continued fraction expansions for series involving known combinatorial quantities the Catalan numbers, the Bell and Stifting numbers, the tangent and secant numbers, the Bill are all Diadrian numbers. We allow the combinatorial interpretations for the coefficients of the elliptic functions, the coefficients of inverses of the Tohebycheff, Chartier, Hermite, Laguerre and Meixner polynomials. Other applications include cycles of binomial coefficients and inversion formulae. Most of the proofs follow from direct generative predistories include vectore objects. We analyse a simple storage allocation scheme in which two stacks grow and shrink inside a shared memory area. To that purpose, we provide analytic expressions for the number of 2-dimensional random walks in a triangle with two reflecting barriers and one absorbing barrier.

We obtain probability distributions and expectations of characteristic parameters of that shared memory scheme, namely the sizes of the stacks and the time until the system runs out of memory.

This provides a complete solution to an open problem posed by Knuth in "The Art of Computer Programming", Vol. 1, 1968 [Ex. 2.2.2.13].



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Theoretical Computer Science 49 (1987) 283-309 North-Holland

# ANALYTIC MODELS AND AMBIGUITY OF CONTEXT-FREE LANGUAGES\*

Philippe FLAJOLET INRIA, Rocquencourt, 78150 Le Chesnay Cedex, France

Abstract. We establish that several classical context-free languages are inherently ambiguous by proving that their counting generating functions, when considered as analytic functions, exhibit some characteristic form of transcendental behaviour. To that purpose, we survey some general results on elementary analytic properties and enumerative uses of algebraic functions in relation to formal languages. In particular, the paper contains a general density theorem for unambiguous context-free languages. FABRICE GUILLEMIN,\*\* France Telecom

THE FORMAL THEORY OF BIRTH-AND-DEATH PROCESSES, LATTICE PATH COMBINATORICS AND CONTINUED FRACTIONS PHILIPPE FLAJOLET,\* INVIA

#### Abstract

Classic works of Karlin and McGregor and Jones and Magnus have established a general correspondence between continuous-time birth and-teath processes and continued fractions of the Stielijes-Jacobi type together with their associated orthogonal polynomials. This finadmental carlies of the state of the state state of the basic relation between weighted lattice paths and continued fractions of herwise known from combinitorial theory. Given that sample paths of the embedded Matrico vhain of a birth-and-eath process are lattice paths. Laplace transforms of a number of transient characteristics can observe the state polynomial. Applications include the analysis of robuttoms in a state state, area, or number of transient of the accentric condition is statisfied.

Keywords: Lattice path combinatorics; continued fractions; orthogonal polynomials; birth-and-death process; first passage time; excursions; transient characteristics.





Theoretical Computer Science

Theoretical Computer Science 281 (2002) 37-80

Basic analytic combinatorics of directed lattice paths

Cyril Banderier\*, Philippe Flajolet Algorithms Project, INRIA, Rocquencourt, 78150 Le Chesnay, France

#### Abstract

This paper develops a unified enumerative and asymptotic theory of directed two-dimensional lattice parks in hisf-planes and quarterplanes. The lattice parks are specified by a finite set of rules that are both time and space homogeneous, and have a privileged direction of increase. (They are then essentially one-dimensional objects.) The theory relies on a specific "kernell method" that provides an important decomposition of the algebraic generating functions involved, precise computed scientification of the mathematication of the state of the specific asymptotic precise computed scientification of the state of a given length ander various accstraints (fordges, excursions, meanders) as well as a characterization of the limit laws associated to several basic parameters of patks. (20 202 Elsevier Science BV. All rights reserved.

Keywords: Lattice path; Analytic combinatorics; Kernel method; Singularity analysis; Generalized ballot problem; Catalan numbers

Séminaire Lotharingien de Combinatoire 54 (2006), Article B54g

#### THE FERMAT CUBIC, ELLIPTIC FUNCTIONS, CONTINUED FRACTIONS, AND A COMBINATORIAL EXCURSION

#### ERIC VAN FOSSEN CONRAD AND PHILIPPE FLAJOLET

Kindly dedicated to Gérard ··· Xavier Viennot on the occasion of his siztieth birthday.

Asymptot: Elliptic functions considered by Dison in the intertexth contrast and the observations of the state of the stat



The Annals of Probability 2005, Vol. 33, No. 3, 1200–1233 DOI 10.1214/009117905000000026 © Institute of Mathematical Statistics, 2005

#### ANALYTIC URNS

#### BY PHILIPPE FLAJOLET, JOAQUIM GABARRÓ AND HELMUT PEKARI

#### INRIA Rocquencourt, Universitat Politècnica de Catalunya and Universitat Politècnica de Catalunya

This article describes a purely analytic approach to urn models of the generalized or extended Pólya-Eggenberger type, in the case of two types of balls and constant "balance," that is, constant row sum. The treatment starts from a quasilinear first-order partial differential equation associated with a combinatorial renormalization of the model and bases itself on elementary conformal mapping arguments coupled with singularity analysis techniques. Probabilistic consequences in the case of "subtractive" urns are new representations for the probability distribution of the urn's composition at any time n, structural information on the shape of moments of all orders. estimates of the speed of convergence to the Gaussian limit and an explicit determination of the associated large deviation function. In the general case, analytic solutions involve Abelian integrals over the Fermat curve  $x^{h} + y^{h} = 1$ . Several urn models, including a classical one associated with balanced trees (2-3 trees and fringe-balanced search trees) and related to a previous study of Panholzer and Prodinger, as well as all urns of balance 1 or 2 and a sporadic urn of balance 3, are shown to admit of explicit representations in terms of Weierstraß elliptic functions: these elliptic models appear precisely to correspond to regular tessellations of the Euclidean plane.

Fourth Colloquium on Mathematics and Computer Science

DMTCS proc. AG, 2006, 59-118

#### Some exactly solvable models of urn process theory

#### Philippe Flajolet, Philippe Dumas, and Vincent Puyhaubert

Algorithms Project, INRIA, F-78153 Le Chesnay (France)

We enablish a fundamental iconception between discrete dime balanced unp processes and certain ordinary differeing large-mass-which are modificant automously and of a single momental form. As a conception, all balanced may proceedings are expressed in terms of certain Arban mitigated in the single sector of the first term. They exhibit an absolute generic functional, or generic dimensional and the single sector of the first sector of the first sector of the single sector of the single



J. Symbolic Computation (1995) 20, 653-671

#### Computer Algebra Libraries for Combinatorial Structures

PHILIPPE FLAJOLET AND BRUNO SALVY

Algorithms Project, INRIA, 78153 Le Chesnay, France

(Received 29 December 1994)

This paper introduces the framework of decompossibe combinatorial structures and their traversal algorithms. A combinisterial type in docomposable if it admits a specification in terms of unions, products, sequences, sets, and cycles, either in the labelled or in the unlabelled cortex. Have, ropersite of decomposable structures are decidable. Generating function equations, counting sequences, and random generation algorithms can be comresonably large schedules. Mogle Hinter task independent such decision procedures are briefly surveysed (LDR, construct, equivalent). In addition, libraries for manipulating holonomic sequences and functions are presented (gran, gran).

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Journal of Symbolic Computation

Journal of Symbolic Computation 41 (2006) 1-29

www.elsevier.com/locate/jsc

#### Fast computation of special resultants

Alin Bostan<sup>a,\*</sup>, Philippe Flajolet<sup>a</sup>, Bruno Salvy<sup>a</sup>, Éric Schost<sup>b</sup>

<sup>8</sup> Algorithms Project, Inria Rocquencourt, 78153 Le Chesnay, France <sup>b</sup> LIX, École polytechnique, 91128 Palaiseau, France

> Received 3 September 2003; accepted 9 July 2005 Available online 25 October 2005

#### Abstract

We propose fast algorithms for computing composed products and composed sums, as well as diamond products of univariate polynomials. These operations correspond to special multivariate resultants, that we compute using power sums of roots of polynomials, by means of their generating series. © 2005 Elsevier Ltd. All rights reserved.

Keywords: Diamond product; Composed product; Composed sum; Complexity; Tellegen's principle

▷ Nearest-neighbor walks in the quarter plane  $\mathbb{N}^2$ ; admissible steps  $\mathfrak{S} \subseteq \{\checkmark, \leftarrow, \nwarrow, \uparrow, \nearrow, \rightarrow, \searrow, \downarrow\}.$ 

 $\triangleright \mathfrak{S}$ -walks = walks in  $\mathbb{N}^2$  starting at (0,0) and using steps in  $\mathfrak{S}$ .

 $\triangleright$  Nearest-neighbor walks in the quarter plane  $\mathbb{N}^2$ ; admissible steps

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▷  $\mathfrak{S}$ -walks = walks in  $\mathbb{N}^2$  starting at (0,0) and using steps in  $\mathfrak{S}$ . ▷  $f_{\mathfrak{S}}(n; i, j)$  = number of  $\mathfrak{S}$ -walks ending at (i, j) and consisting of exactly *n* steps.

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$$F_{\mathfrak{S}}(t;x,y) = \sum_{n=0}^{\infty} \left( \sum_{i,j=0}^{\infty} f_{\mathfrak{S}}(n;i,j) x^{i} y^{j} \right) t^{n} \in \mathbb{Q}[x,y][[t]]$$

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**Questions:** Given  $\mathfrak{S}$ , what can be said about  $F_{\mathfrak{S}}(t; x, y)$ ? Structure? (algebraic / holonomic) Explicit form? Asymptotics?

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 $F_{\mathfrak{S}}(t; 0, 0) \sim \text{counts } \mathfrak{S}\text{-walks returning to the origin (excursions)};$  $F_{\mathfrak{S}}(t; 1, 1) \sim \text{counts } \mathfrak{S}\text{-walks with prescribed length};$  $F_{\mathfrak{S}}(t; 1, 0) \sim \text{counts } \mathfrak{S}\text{-walks ending on the horizontal axis.}$ 



There are  $2^8$  such sets.



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Some of these  $2^8 = 256$  step sets are:



trivial,



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trivial, simple,



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Finally, there remain 79 inherently different cases!

## Two important cases: Kreweras and Gessel walks

$$\mathfrak{S} = \{\downarrow, \leftarrow, \nearrow\} \qquad F_{\mathfrak{S}}(t; x, y) \equiv K(t; x, y)$$

$$\mathfrak{S} = \{\nearrow, \checkmark, \leftarrow, \rightarrow\} \quad F_{\mathfrak{S}}(t; x, y) \equiv \mathsf{G}(t; x, y)$$





Example: A Kreweras excursion.





*Holonomic*:  $S(t) \in \mathbb{Q}[[t]]$  satisfying a linear differential equation with polynomial coefficients  $c_r(t)S^{(r)}(t) + \cdots + c_0(t)S(t) = 0$ .



Holonomic:  $S(t) \in \mathbb{Q}[[t]]$  satisfying a linear differential equation with polynomial coefficients  $c_r(t)S^{(r)}(t) + \cdots + c_0(t)S(t) = 0$ . Algebraic:  $S(t) \in \mathbb{Q}[[t]]$  root of a polynomial  $P \in \mathbb{Q}[t, T]$ .



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$$_{2}F_{1}\begin{pmatrix} a & b \\ c \end{pmatrix} t = \sum_{n=0}^{\infty} \frac{(a)_{n}(b)_{n}}{(c)_{n}} \frac{t^{n}}{n!}, \quad (a)_{n} = a(a+1)\cdots(a+n-1).$$



 $S \in \mathbb{Q}[[x, y, t]]$  is *holonomic* if the set of all partial derivatives of S spans a finite-dimensional vector space over  $\mathbb{Q}(x, y, t)$ .



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 $S \in \mathbb{Q}[[x, y, t]]$  is *algebraic* if it is the root of a  $P \in \mathbb{Q}[x, y, t, T]$ .

Theorem [Kreweras 1965; 100 pages combinatorial proof!]  

$$K(t; 0, 0) = {}_{3}F_{2} \begin{pmatrix} 1/3 & 2/3 & 1 \\ 3/2 & 2 \end{pmatrix} | 27 t^{3} \end{pmatrix} = \sum_{n=0}^{\infty} \frac{4^{n} \binom{3n}{n}}{(n+1)(2n+1)} t^{3n}.$$

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Theorem [Gessel's conjecture; Kauers, Koutschan, Zeilberger 2009]  $G(t; 0, 0) = {}_{3}F_{2} \begin{pmatrix} 5/6 & 1/2 & 1 \\ 5/3 & 2 \end{pmatrix} | 16t^{2} \end{pmatrix} = \sum_{n=0}^{\infty} \frac{(5/6)_{n}(1/2)_{n}}{(5/3)_{n}(2)_{n}} (4t)^{2n}.$ 

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► Computer-driven discovery and proof; no human proof yet.

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► G(t; x, y) had been conjectured to be non-holonomic.
Theorem [B. & Kauers'10] G(t; x, y) is holonomic, even algebraic.
► Fresh news: human proof just announced [B., Kurkova, Raschel].

# Main results (II): Explicit form for G(t; x, y)

Theorem Let 
$$V = 1 + 4t^2 + 36t^4 + 396t^6 + \cdots$$
 be a root of  
 $(V - 1)(1 + 3/V)^3 = (16t)^2$ ,  
let  $U = 1 + 2t^2 + 16t^4 + 2xt^5 + 2(x^2 + 83)t^6 + \cdots$  be a root of  
 $x(V - 1)(V + 1)U^3 - 2V(3x + 5xV - 8Vt)U^2$   
 $-xV(V^2 - 24V - 9)U + 2V^2(xV - 9x - 8Vt) = 0$ ,  
let  $W = t^2 + (y + 8)t^4 + 2(y^2 + 8y + 41)t^6 + \cdots$  be a root of  
 $y(1 - V)W^3 + y(V + 3)W^2 - (V + 3)W + V - 1 = 0$ .  
Then  $G(t; x, y)$  is equal to

$$\frac{\frac{64(U(V+1)-2V)V^{3/2}}{x(U^2-V(U^2-8U+9-V))^2} - \frac{y(W-1)^4(1-Wy)V^{-3/2}}{t(y+1)(1-W)(W^2y+1)^2}}{(1+y+x^2y+x^2y^2)t-xy} - \frac{1}{tx(y+1)}$$

▷ Computer-driven discovery and proof; no human proof yet.

# Main results (II): Explicit form for G(t; 0, 0)

Theorem Let  $V = 1 + 4t^2 + 36t^4 + 396t^6 + \cdots$  be a root of  $(V - 1)(1 + 3/V)^3 = (16t)^2$ . Then G(t; 0, 0) is equal to  $\frac{\frac{V^2 + 6V - 3}{4V^{3/2}} - 1}{2t^2}$ .
### Main results (II): Explicit form for G(t; 0, 0)

Theorem Let  $V = 1 + 4t^2 + 36t^4 + 396t^6 + \cdots$  be a root of  $(V - 1)(1 + 3/V)^3 = (16t)^2$ . Then G(t; 0, 0) is equal to  $\frac{\frac{V^2 + 6V - 3}{4V^{3/2}} - 1}{2t^2}$ .

De : Philippe Flajolet Objet : Rép : ca y est... Date : 7 août 2008 09:49:21 HAEC À : Alin Bostan Cc : Bruno Salvy . Philippe Flajolet

La courbe P(x,y)=0 est de genre 0 et possede une merveilleuse parametrisation rationnelle:

 $x = -(t-1)^{*}(t^{2}-t+1)^{3}/(t^{2}^{*}(t-2)^{6}), y = t^{*}(t^{3}-2)^{*}(t-2)^{3}/((t-1)^{*}(t^{2}-t+1)^{3});$ 

Y a de la structure....

# 

| OEIS Tag | Sample step set | Equation sizes |       |      | OEIS Tag | Sample step set | Eq    | es    |      |
|----------|-----------------|----------------|-------|------|----------|-----------------|-------|-------|------|
| A000012  | · · •           | 1, 0           | 1, 1  | 1, 1 | A000079  | - ••            | 1, 0  | 1, 1  | 1, 1 |
| A001405  | •••             | 2, 1           | 2, 3  | 2, 2 | A000244  |                 | 1, 0  | 1, 1  | 1, 1 |
| A001006  | •               | 2, 1           | 2, 3  | 2, 2 | A005773  | •••             | 2, 1  | 2, 3  | 2, 2 |
| A126087  | •••<br>••       | 3, 1           | 2, 5  | 2, 2 | A151255  | • · · ·         | 6, 8  | 4, 16 | -    |
| A151265  | 1.1<br>1.1      | 6, 4           | 4, 9  | 6, 8 | A151266  | •               | 7,10  | 5, 16 | -    |
| A151278  |                 | 7,4            | 4, 12 | 6, 8 | A151281  | •••             | 3, 1  | 2, 5  | 2, 2 |
| A005558  |                 | 2, 3           | 3, 5  | -    | A005566  |                 | 2, 2  | 3, 4  | -    |
| A018224  | •••<br>•••      | 2, 3           | 3, 5  | -    | A060899  |                 | 2, 1  | 2, 3  | 2, 2 |
| A060900  |                 | 2, 3           | 3, 5  | 8, 9 | A128386  | <b>.</b>        | 3, 1  | 2, 5  | 2, 2 |
| A129637  | ••••<br>• • •   | 3, 1           | 2, 5  | 2, 2 | A151261  |                 | 5,8   | 4, 15 | -    |
| A151282  | ••••            | 3, 1           | 2, 5  | 2, 2 | A151291  |                 | 6, 10 | 5, 15 | -    |
| A151275  |                 | 9, 18          | 5,24  | -    | A151287  | ::              | 7, 11 | 5,19  | -    |
| A151292  | <b>.</b>        | 3, 1           | 2, 5  | 2, 2 | A151302  | •••             | 9,18  | 5, 24 | -    |
| A151307  |                 | 8, 15          | 5,20  | -    | A151318  |                 | 2, 1  | 2, 3  | 2, 2 |
| A129400  | •               | 2, 1           | 2, 3  | 2, 2 | A151297  |                 | 7, 11 | 5, 18 | -    |
| A151312  |                 | 4, 5           | 3, 8  | -    | A151323  |                 | 2, 1  | 2, 3  | 4, 4 |
| A151326  |                 | 7,14           | 5,18  | -    | A151314  |                 | 9, 18 | 5, 24 | -    |
| A151329  | ::              | 9, 18          | 5,24  | -    | A151331  |                 | 3, 4  | 3, 6  | -    |

Equation sizes = {order, degree}(rec, diffeq, algeq).

▷ Computer-driven; confirmed by human proofs in [Bousquet-Mélou & Mishna, 2010].

| ~        |             |       | 455       |      |   |          |             |       |                |      |   |  |
|----------|-------------|-------|-----------|------|---|----------|-------------|-------|----------------|------|---|--|
| OEIS Tag | Steps       | Equ   | ation siz | es   | Asymptotics   | OEIS Tag | Tag Steps   |       | Equation sizes |      | Asymptotics   |  |
| A000012  | :: <b>:</b> | 1, 0  | 1, 1      | 1, 1 | 1   | A000079  | : <b>••</b> | 1,0   | 1, 1           | 1, 1 | $2^n$   |  |
| A001405  | :::         | 2, 1  | 2, 3      | 2, 2 | $\frac{\sqrt{2}}{\Gamma(\frac{1}{2})}\frac{2^n}{\sqrt{n}}$              | A000244  | :::         | 1,0   | 1, 1           | 1, 1 | $3^n$   |  |
| A001006  | •:          | 2, 1  | 2, 3      | 2, 2 | $\frac{3\sqrt{3}}{2\Gamma(\frac{1}{2})}\frac{3^n}{n^{3/2}}$             | A005773  |             | 2, 1  | 2, 3           | 2, 2 | $\frac{\sqrt{3}}{\Gamma(\frac{1}{2})}\frac{3^n}{\sqrt{n}}$          |  |
| A126087  | •           | 3, 1  | 2, 5      | 2, 2 | $\frac{12\sqrt{2}}{\Gamma(\frac{1}{2})}\frac{2^{3n/2}}{n^{3/2}}$        | A151255  | ::          | 6, 8  | 4, 16          | -    | $\frac{24\sqrt{2}}{\pi} \frac{2^{3n/2}}{n^2}$                       |  |
| A151265  | •••         | 6, 4  | 4, 9      | 6, 8 | $\frac{2\sqrt{2}}{\Gamma(\frac{1}{4})} \frac{3^n}{n^{3/4}}$             | A151266  | •::         | 7, 10 | 5, 16          | -    | $\frac{\sqrt{3}}{2\Gamma(\frac{1}{2})}\frac{3^n}{\sqrt{n}}$         |  |
| A151278  |             | 7, 4  | 4, 12     | 6, 8 | $\frac{3\sqrt{3}}{\sqrt{2}\Gamma(\frac{1}{4})}\frac{3^{n}}{n^{3/4}}$    | A151281  | :::         | 3, 1  | 2, 5           | 2, 2 | $\frac{1}{2}3^{n}$  |  |
| A005558  | •           | 2, 3  | 3, 5      | -    | $\frac{8}{\pi} \frac{4^n}{n^2}$   | A005566  |             | 2, 2  | 3, 4           | -    | $\frac{4}{\pi} \frac{4^n}{n}$                                       |  |
| A018224  | ::          | 2, 3  | 3, 5      | -    | $\frac{2}{\pi} \frac{4^n}{n}$   | A060899  | ::          | 2, 1  | 2, 3           | 2, 2 | $\frac{\sqrt{2}}{\Gamma(\frac{1}{2})} \frac{4^n}{\sqrt{n}}$         |  |
| A060900  | ::          | 2, 3  | 3, 5      | 8, 9 | $\frac{4\sqrt{3}}{3\Gamma(\frac{1}{3})}\frac{4^{n}}{n^{2/3}}$           | A128386  | ::          | 3, 1  | 2, 5           | 2, 2 | $\frac{6\sqrt{2}}{\Gamma(\frac{1}{2})} \frac{2^n 3^{n/2}}{n^{3/2}}$ |  |
| A129637  | :::         | 3, 1  | 2, 5      | 2, 2 | $\frac{1}{2}4^{n}$  | A151261  | ::          | 5,8   | 4, 15          | -    | $\frac{12\sqrt{3}}{\pi} \frac{2^n 3^{n/2}}{n^2}$                    |  |
| A151282  | ••••        | 3, 1  | 2, 5      | 2, 2 | $\frac{A^2 B^{3/2}}{2^{3/4} \Gamma(\frac{1}{2})} \frac{B^n}{n^{3/2}}$   | A151291  | •::         | 6, 10 | 5, 15          | -    | $\frac{4}{3\Gamma(\frac{1}{2})}\frac{4^n}{\sqrt{n}}$                |  |
| A151275  | ::          | 9, 18 | 5, 24     | -    | $\frac{12\sqrt{30}}{\pi} \frac{(\sqrt{24})^n}{n^2}$                     | A151287  | ::          | 7, 11 | 5, 19          | -    | $\frac{\sqrt{8}A^{7/2}}{\pi} \frac{(2A)^n}{n^2}$                    |  |
| A151292  | :::         | 3, 1  | 2, 5      | 2, 2 | $\frac{\sqrt[4]{3}C^2D^{3/2}}{8\Gamma(\frac{1}{2})}\frac{D^n}{n^{3/2}}$ | A151302  |             | 9, 18 | 5, 24          | -    | $\frac{\sqrt{5}}{3\sqrt{2}\Gamma(\frac{1}{2})}\frac{5^n}{\sqrt{n}}$ |  |
| A151307  | •           | 8, 15 | 5, 20     | -    | $\frac{\sqrt{5}}{2\sqrt{2}\Gamma(\frac{1}{2})}\frac{5^n}{\sqrt{n}}$     | A151318  |             | 2, 1  | 2, 3           | 2, 2 | $\frac{\sqrt{5/2}}{\Gamma(\frac{1}{2})} \frac{5^n}{\sqrt{n}}$       |  |
| A129400  | •::         | 2, 1  | 2, 3      | 2, 2 | $\frac{3\sqrt{3}}{2\Gamma(\frac{1}{2})} \frac{6^n}{n^{3/2}}$            | A151297  |             | 7, 11 | 5, 18          | -    | $\frac{\sqrt{3}C^{7/2}}{2\pi} \frac{(2C)^n}{n^2}$                   |  |
| A151312  | ::          | 4, 5  | 3, 8      | -    | $\frac{\sqrt{6}}{\pi} \frac{6^n}{n}$                                    | A151323  | ::*         | 2, 1  | 2, 3           | 4,4  | $\frac{\sqrt{2} 3^{3/4}}{\Gamma(\frac{1}{4})} \frac{6^n}{n^{3/4}}$  |  |
| A151326  | ::::        | 7, 14 | 5, 18     | -    | $\frac{2\sqrt{3}}{3\Gamma(\frac{1}{2})}\frac{6^n}{\sqrt{n}}$            | A151314  |             | 9, 18 | 5, 24          | -    | $\frac{EF^{7/2}}{5\sqrt{95}\pi} \frac{(2F)^n}{n^2}$                 |  |
| A151329  |             | 9, 18 | 5, 24     | -    | $\frac{\sqrt{7/3}}{3\Gamma(\frac{1}{2})}\frac{7^n}{\sqrt{n}}$           | A151331  |             | 3,4   | 3, 6           | -    | $\frac{8}{3\pi} \frac{8^n}{n}$                                      |  |

#### Experimental classification of walks with holonomic $F_{\mathfrak{S}}(t; 1, 1)$

▷ Computer-driven; recent human proofs of asymptotics by [Fayolle & Raschel, 2012]. 16/54

### The group of a walk: an example



The characteristic polynomial  $\chi_{\mathfrak{S}} = x + \frac{1}{x} + y + \frac{1}{y}$ 

### The group of a walk: an example



The characteristic polynomial  $\chi_{\mathfrak{S}} = x + \frac{1}{x} + y + \frac{1}{y}$  is left invariant under

$$\psi(x,y) = \left(x, \frac{1}{y}\right), \qquad \phi(x,y) = \left(\frac{1}{x}, y\right),$$

### The group of a walk: an example



The characteristic polynomial  $\chi_{\mathfrak{S}} = x + \frac{1}{x} + y + \frac{1}{y}$  is left invariant under

$$\psi(x,y) = \left(x, \frac{1}{y}\right), \qquad \phi(x,y) = \left(\frac{1}{x}, y\right),$$

and thus under any element of the group

$$\langle \psi, \phi \rangle = \left\{ (x, y), \left( x, \frac{1}{y} \right), \left( \frac{1}{x}, \frac{1}{y} \right), \left( \frac{1}{x}, y \right) \right\}.$$

### The group of a walk: the general case



The polynomial  $\chi_{\mathfrak{S}} := \sum_{(i,j)\in\mathfrak{S}} x^i y^j = \sum_{i=-1}^1 B_i(y) x^i = \sum_{j=-1}^1 A_j(x) y^j$ 

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is left invariant under

$$\psi(x,y) = \left(x, \frac{A_{-1}(x)}{A_{+1}(x)}\frac{1}{y}\right), \qquad \phi(x,y) = \left(\frac{B_{-1}(y)}{B_{+1}(y)}\frac{1}{x}, y\right),$$

#### The group of a walk: the general case



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and thus under any element of the group

$$\mathcal{G}_{\mathfrak{S}} := \langle \psi, \phi \rangle.$$



Order 4,



Order 4, order 6,



Order 4,

order 6,

order 8,



79 step sets







## The 23 cases with a finite group

(i) 16 with a vertical symmetry, and group isomorphic to D<sub>2</sub>
If with a vertical symmetry, and group isomorphic to D<sub>2</sub>
If with a vertical symmetry, and group isomorphic to D<sub>2</sub>
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If with a vertical symm

(ii) 5 with a *diagonal* or an *anti-diagonal symmetry*, and group isomorphic to  $D_3$ 



(iii) 2 with group isomorphic to  $D_4$ 



In red, cases with algebraic generating functions (ii)+(iii): zero drift  $\sum_{s \in G} s$ 

# Main results (IV): explicit expressions for the 19 holonomic transcendental cases

#### Theorem [B.-Chyzak-Van Hoeij-Kauers-Pech, 2011]

Let  $\mathfrak{S}$  be one of the 19 step sets with finite group  $\mathcal{G}_{\mathfrak{S}}$ , and such that the generating series  $F = F_{\mathfrak{S}}(t; x, y)$  is not algebraic. Then F is expressible using iterated integrals of  $_2F_1$  expressions.

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*Example* (King walks in the quarter plane, A025595)

$$F_{\text{res}}(t;1,1) = \frac{1}{t} \int_0^t \frac{1}{(1+4x)^3} \cdot {}_2F_1\left(\frac{3}{2},\frac{3}{2} \mid \frac{16x(1+x)}{(1+4x)^2}\right) dx$$

 $= 1 + 3t + 18t^{2} + 105t^{3} + 684t^{4} + 4550t^{5} + 31340t^{6} + 219555t^{7} + \cdots$ 

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 $= 1 + 3t + 18t^{2} + 105t^{3} + 684t^{4} + 4550t^{5} + 31340t^{6} + 219555t^{7} + \cdots$ 

Computer-driven discovery and proof; no human proof yet.
 Proof uses creative telescoping, ODE factorization, ODE solving.

# Main results (V): algorithmic proof of non-holonomy for the 51 non-singular cases with infinite group

### Theorem [B.-Rachel-Salvy, 2012]

Let  $\mathfrak{S}$  be one of the 51 step sets with infinite group  $\mathcal{G}_{\mathfrak{S}}$ , and such that the excursions series  $F_{\mathfrak{S}}(t;0,0)$  is not equal to 1. Then  $F_{\mathfrak{S}}(t;0,0)$ , and in particular  $F_{\mathfrak{S}}(t;x,y)$ , are non-holonomic.

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 Algorithmic, computer-driven, proof. Uses Gröbner basis computations, polynomial factorization, cyclotomy testing.
 Based on *two ingredients*: asymptotics + irrationality.

▷ [Kurkova & Raschel 2012] Alternative proof of  $F_{\mathfrak{S}}(t; x, y)$  is non-holonomic. No human proof yet for  $F_{\mathfrak{S}}(t; 0, 0)$  non-holonomic.

### The 56 cases with infinite group



In blue, non-singular cases, solved by [B., Raschel & Salvy, 2012] In red, singular cases, solved by [Melczer & Mishna 2012]

### Summary – classification of 2D non-singular walks

The Big Theorem Let  $\mathfrak{S}$  be one of the 74 non-singular step sets. The following assertions are equivalent:

- (1) The full generating series  $F_{\mathfrak{S}}(t; x, y)$  is holonomic
- (2) the excursions generating series  $F_{\mathfrak{S}}(t; 0, 0)$  is holonomic
- (3) the excursions seq.  $[t^n] F_{\mathfrak{S}}(t; 0, 0)$  is  $\sim K \cdot \rho^n \cdot n^{\alpha}$ , with  $\alpha \in \mathbb{Q}$
- (4) the group  $\mathcal{G}_{\mathfrak{S}}$  is finite (and  $|\mathcal{G}_{\mathfrak{S}}| = 2 \cdot \min\{\ell \in \mathbb{N}^* \mid \frac{\ell}{\alpha+1} \in \mathbb{Z}\})$
- (5) the step set  $\mathfrak{S}$  has either an axial symmetry, or a zero drift

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Moreover, under (1)–(5),  $F_{\mathfrak{S}}(t; x, y)$  is algebraic if and only if the step set  $\mathfrak{S}$  has positive covariance  $\sum_{(i,j)\in\mathfrak{S}} ij - \sum_{(i,j)\in\mathfrak{S}} i \cdot \sum_{(i,j)\in\mathfrak{S}} j > 0$ 

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In this case,  $F_{\mathfrak{S}}(t; x, y)$  is expressible using nested radicals. If not,  $F_{\mathfrak{S}}(t; x, y)$  is expressible using iterated integrals of  $_2F_1$  expressions.

## Main methods

### (1) for proving non-holonomy

- (1a) Infinite number of singularities, or lacunary
- (1b) Asymptotics

### (2) for proving holonomy

- (2a) Diagonals, or positive parts, of rational functions
- (2b) Guess'n'Prove

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### (1) for proving non-holonomy

- (1a) Infinite number of singularities, or lacunary
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### (2) for proving holonomy

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- (2b) Guess'n'Prove

▷ All methods are algorithmic.

# Methodology for proving algebraicity

### Methodology for proving algebraicity

Experimental mathematics -Guess'n'Prove- approach:

- (S1) high order expansion of the generating series F<sub>G</sub>(t; x, y);
  (S2) guessing candidates for minimal polynomials of F<sub>G</sub>(t; x, 0) and F<sub>G</sub>(t; 0, y), based on Hermite-Padé approximation;
- (S3) rigorous certification of the minimal polynomials, based on (exact) polynomial computations.

# Step (S1): high order series expansions

 $f_{\mathfrak{S}}(n; i, j)$  satisfies the recurrence with constant coefficients  $f_{\mathfrak{S}}(n+1; i, j) = \sum_{(u,v)\in\mathfrak{S}} f_{\mathfrak{S}}(n; i-u, j-v) \text{ for } n, i, j \ge 0$ 

+ init. cond.  $f_{\mathfrak{S}}(0; i, j) = \delta_{0, i, j}$  and  $f_{\mathfrak{S}}(n; -1, j) = f_{\mathfrak{S}}(n; i, -1) = 0$ .

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$$k(n+1; i, j) = k(n; i+1, j) + k(n; i, j+1) + k(n; i-1, j-1)$$



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▷ Recurrence is used to compute  $F_{\mathfrak{S}}(t; x, y) \mod t^N$  for large N.

$$\begin{split} \mathcal{K}(t;x,y) &= 1 + xyt + (x^2y^2 + y + x)t^2 + (x^3y^3 + 2xy^2 + 2x^2y + 2)t^3 \\ &+ (x^4y^4 + 3x^2y^3 + 3x^3y^2 + 2y^2 + 6xy + 2x^2)t^4 \\ &+ (x^5y^5 + 4x^3y^4 + 4x^4y^3 + 5xy^3 + 12x^2y^2 + 5x^3y + 8y + 8x)t^5 + \cdots \end{split}$$

# Step (S2): guessing equations for $F_{\mathfrak{S}}(t; x, y)$ , a first idea

In terms of generating series, the recurrence on k(n; i, j) reads

$$(xy - (x + y + x^2y^2)t)K(t; x, y)$$
  
=  $xy - xt K(t; x, 0) - yt K(t; 0, y)$  (KerEq)

▷ This *kernel equation* can be seen as a multivariate analogue of  $(1 - t - t^2) \cdot \sum_{n \ge 0} \ell_n t^n = 1$ , where  $\ell_n$  are the Fibonacci numbers.

▷ A similar kernel equation holds for  $F_{\mathfrak{S}}(t; x, y)$ , for any  $\mathfrak{S}$ -walk.

*Corollary.*  $F_{\mathfrak{S}}(t; x, y)$  is holonomic (resp. algebraic) if and only if  $F_{\mathfrak{S}}(t; x, 0)$  and  $F_{\mathfrak{S}}(t; 0, y)$  are both holonomic (resp. algebraic).

▷ Crucial simplification: equations for G(t; x, y) are huge ( $\approx$  30Gb)

# **Step (S2): guessing equations for** $F_{\mathfrak{S}}(t; x, 0) \& F_{\mathfrak{S}}(t; 0, y)$

*Task 1:* Given the first *N* terms of  $S = F_{\mathfrak{S}}(t; x, 0) \in \mathbb{Q}[x][[t]]$ , search for a *differential equation* satisfied by *S* at precision *N*:

$$\mathcal{L}_{x,0}(S) = c_r(x,t) \cdot \frac{\partial^r S}{\partial t^r} + \dots + c_1(x,t) \cdot \frac{\partial S}{\partial t} + c_0(x,t) \cdot S = 0 \mod t^N$$
## **Step (S2): guessing equations for** $F_{\mathfrak{S}}(t; x, 0) \& F_{\mathfrak{S}}(t; 0, y)$

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- Both tasks amount to linear algebra in size N over  $\mathbb{Q}(x)$ .
- In practice, we use modular Hermite-Padé approximation (Beckermann-Labahn algorithm) combined with (rational) evaluation-interpolation and rational number reconstruction.
- ▶ Fast (FFT-based) arithmetic in  $\mathbb{F}_{\rho}[t]$  and  $\mathbb{F}_{\rho}[t]\langle \frac{t}{\partial t}\rangle$ .

## **Step (S2): guessing equations for** G(t; x, 0) and G(t; 0, y)

Using N = 1200 terms of G(t; x, y), we guessed candidates

- ▶  $\mathcal{P}_{x,0}$  in  $\mathbb{Z}[x, t, T]$  of tridegree (32, 43, 24), 21 digits coefficients
- ▶  $\mathcal{P}_{0,y}$  in  $\mathbb{Z}[y, t, T]$  of tridegree (40, 44, 24), 23 digits coefficients

such that

 $\mathcal{P}_{x,0}(x, t, G(t; x, 0)) = \mathcal{P}_{0,y}(y, t, G(t; 0, y)) = 0 \bmod t^{1200}.$ 

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▷ We actually first guessed *differential equations*<sup>†</sup>, then computed their *p*-*curvatures* to empirically certify them. This led to suspect the algebraicity of G(t; x, 0) and G(t; 0, y), using a conjecture of Grothendieck (on differential equations modulo *p*) as an oracle.

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▷ Guessing  $\mathcal{P}_{x,0}$  by *undetermined coefficients* would require solving a dense linear system of size  $\approx 100\,000$ , and  $\approx 1000$  digits entries!

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 Implicit function theorem: ∃! root r(t) ∈ Q[[t]] of P.
 r(t)=∑<sub>n=0</sub><sup>∞</sup> r<sub>n</sub>t<sup>n</sup> being algebraic, it is holonomic, and so is (r<sub>n</sub>):

$$(n+2)(3n+5)r_{n+1}-4(6n+5)(2n+1)r_n=0, \qquad r_0=1$$

 $\Rightarrow$  solution  $r_n = \frac{(5/6)_n(1/2)_n}{(5/3)_n(2)_n} 4^{2n} = g_n$ , thus g(t) = r(t) is algebraic.

1. Setting  $y_0 = \frac{x - t - \sqrt{x^2 - 2tx + t^2(1 - 4x^3)}}{2tx^2} = t + \frac{1}{x}t^2 + \frac{x^3 + 1}{x^2}t^3 + \cdots$ in the kernel equation

$$\underbrace{(xy - (x + y + x^2y^2)t)}_{= 0} K(t; x, y) = xy - xtK(t; x, 0) - ytK(t; 0, y)$$

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shows that U = K(t; x, 0) satisfies the *reduced kernel equation* 

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- 4. Resultant computations + verification of initial terms  $\implies U = H(t, x)$  also satisfies (RKerEq).
- 5. Uniqueness:  $H(t, x) = K(t; x, 0) \implies K(t; x, 0)$  is algebraic!

#### Algebraicity of Kreweras walks: our Maple proof in a nutshell

```
[bostan@inria ~]$ maple
    1\^/1
             Maple 17 (APPLE UNIVERSAL OSX)
._|\| |/|_. Copyright (c) Maplesoft, a division of Waterloo Maple Inc. 2013
\ MAPLE / All rights reserved. Maple is a trademark of
 <____> Waterloo Maple Inc.
             Type ? for help.
# HIGH ORDER EXPANSION (S1)
> st,bu:=time(),kernelopts(bytesused):
> f:=proc(n,i,j)
 option remember;
   if i<0 or i<0 or n<0 then 0
    elif n=0 then if i=0 and i=0 then 1 else 0 fi
    else f(n-1,i-1,j-1)+f(n-1,i,j+1)+f(n-1,i+1,j) fi
 end:
> S:=series(add(add(f(k,i,0)*x^i,i=0..k)*t^k,k=0..80),t,80);
# GUESSING (S2)
> libname:=".".libname:gfun:-version();
                                      3.62
> gfun:-seriestoalgeq(S,Fx(t)):
> P:=collect(numer(subs(Fx(t)=T,%[1])),T);
# RIGOROUS PROOF (S3)
> ker := (T,t,x) -> (x+T+x^2*T^2)*t-x*T:
> pol := unapplv(P.T.t.x):
> p1 := resultant(pol(z-T,t,x),ker(t*z,t,x),z):
> p2 := subs(T=x*T,resultant(numer(pol(T/z,t,z)),ker(z,t,x),z)):
> normal(primpart(p1.T)/primpart(p2.T));
                                        1
# time (in sec) and memory consumption (in Mb)
> trunc(time()-st),trunc((kernelopts(bytesused)-bu)/1000^2);
                                     7.618
```

# Step (S3): rigorous proof for Gessel walks

Same philosophy, but several complications:

- ▶ stepset diagonal symmetry is lost:  $G(t; x, y) \neq G(t; y, x)$ ;
- G(t; 0, 0) occurs in (KerEq);
- equations are  $\approx 5\,000$  times bigger.
- $\rightarrow$  replace (RKerEq) by a *system* of 2 reduced kernel equations.

 $\rightarrow$  fast algorithms needed (e.g., [B.-Flajolet-Salvy-Schost'06] for computations with algebraic series).



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#### Fast computation of special resultants

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> Received 3 September 2003; accepted 9 July 2005 Available online 25 October 2005

# Holonomy via the finite group



The polynomial  $\chi_{\mathfrak{S}} = \sum_{(i,j)\in\mathfrak{S}} x^i y^j = x + \frac{1}{x} + y + \frac{1}{y}$ is left invariant under  $(x, y), (\frac{1}{x}, y), (\frac{1}{x}, \frac{1}{y}), (x, \frac{1}{y})$ . The same holds for  $J(t; x, y) = \sum_{(i,j)\in\mathfrak{S}} x^i y^j - 1/t$ .



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J(t; x, y) x y tF(t; x, y) = t x F(t; x, 0) + t y F(t; 0, y) - x y $-J(t; x, y) \frac{1}{x} y tF(t; \frac{1}{x}, y) = -t \frac{1}{x} F(t; \frac{1}{x}, 0) - t y F(t; 0, y) + \frac{1}{x} y$  $J(t; x, y) \frac{1}{x} \frac{1}{y} tF(t; \frac{1}{x}, \frac{1}{y}) = t \frac{1}{x} F(t; \frac{1}{x}, 0) + t \frac{1}{y} F(t; 0, \frac{1}{y}) - \frac{1}{x} \frac{1}{y}$  $-J(t; x, y) x \frac{1}{y} tF(t; x, \frac{1}{y}) = -t x F(t; x, 0) - t \frac{1}{y} F(t; 0, \frac{1}{y}) + x \frac{1}{y}$ 

$$\sum_{\theta \in \mathcal{G}_{\mathfrak{S}}} (-1)^{\theta} \theta \left[ xyt F(t; x, y) \right] =$$

$$\frac{-xy + \frac{1}{x}y - \frac{1}{x}\frac{1}{y} + x\frac{1}{y}}{J(t; x, y)}$$



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is left invariant under  $(x, y), (\frac{1}{x}, y), (\frac{1}{x}, \frac{1}{y}), (x, \frac{1}{y}).$   
The same holds for  $J(t; x, y) = \sum_{(i,j)\in\mathfrak{S}} x^i y^j - 1/t.$ 

$$J(t; x, y) x y tF(t; x, y) = t x F(t; x, 0) + t y F(t; 0, y) - x y$$
  
$$-J(t; x, y) \frac{1}{x} y tF(t; \frac{1}{x}, y) = -t \frac{1}{x} F(t; \frac{1}{x}, 0) - t y F(t; 0, y) + \frac{1}{x} y$$
  
$$J(t; x, y) \frac{1}{x} \frac{1}{y} tF(t; \frac{1}{x}, \frac{1}{y}) = t \frac{1}{x} F(t; \frac{1}{x}, 0) + t \frac{1}{y} F(t; 0, \frac{1}{y}) - \frac{1}{x} \frac{1}{y}$$
  
$$-J(t; x, y) x \frac{1}{y} tF(t; x, \frac{1}{y}) = -t x F(t; x, 0) - t \frac{1}{y} F(t; 0, \frac{1}{y}) + x \frac{1}{y}$$

$$xyt F(t; x, y) = [x^{>}][y^{>}] \frac{-xy + \frac{1}{x}y - \frac{1}{x}\frac{1}{y} + x\frac{1}{y}}{J(t; x, y)}$$

## The 19 transcendental holonomic cases

Theorem [Bousquet-Mélou & Mishna, 2009]

Let  $\mathfrak{S}$  be one of the 19 step sets with finite group  $\mathcal{G}_{\mathfrak{S}}$ , and such that the generating series  $F = F_{\mathfrak{S}}(t; x, y)$  is not algebraic. Then:

$$xyt F(t; x, y) = [x^{>}][y^{>}] \frac{\sum_{\theta \in \mathcal{G}_{\mathfrak{S}}} (-1)^{\theta} \cdot \theta(xy)}{J(t; x, y)}$$

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*Proof*: Use [Lipshitz'88] for positive parts of holonomic series.

If F(t; x, y), as a formal series in t, has its coefficients in  $\mathbb{Q}(x)[y, \frac{1}{y}]$ , then  $[y^{>}]F(t; x, y)$  is algebraic. If in addition  $[y^{>}]F(t; x, y)$  as a formal series in t, has its coefficients in  $\mathbb{Q}[x, \frac{1}{x}, y]$ , then  $[x^{>}][y^{>}]F(t; x, y)$  is holonomic.

#### ▷ Constructive proof, but it leads to a highly inefficient algorithm.

## Explicit expressions for the 19 holonomic cases

#### Theorem [B.-Chyzak-Van Hoeij-Kauers-Pech, 2011]

Let  $\mathfrak{S}$  be one of the 19 step sets with finite group  $\mathcal{G}_{\mathfrak{S}}$ , and such that the generating series  $F = F_{\mathfrak{S}}(t; x, y)$  is not algebraic. Then F is expressible using iterated integrals of  $_2F_1$  expressions.

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#### Sketch of the approach

- 1. (BM&M) If  $R = \sum_{\theta} \frac{(-1)^{\theta} \theta(xy)}{J(t;x,y)}$ , then  $F = \operatorname{Res}_u(\operatorname{Res}_v H)$ , for  $H = \frac{R(t;1/u,1/v)}{(1-xu)(1-yv)}$ .
- 2. If  $P \in \mathbb{Q}(x, y)[t] \langle \partial_t \rangle$  and  $U, V \in \mathbb{Q}(x, y, u, v, t)$  such that  $L(H) = \partial_u U + \partial_v V$ , then  $L(F(t; x, y)) = 0 \longrightarrow$  Chyzak's creative telescoping for finding L.
- 3. Factor L as  $L^{(2)} \cdot L_1^{(1)} \cdots L_t^{(1)}$ , then solve  $L^{(2)}$  in terms of  ${}_2F_1$ s, and deduce F.

# Proofs of non-holonomy

#### A historical example



Theorem [Bousquet-Mélou & Petkovsek'03] The generating series  $F(x, y) = \sum_{m,n \ge 0} a_{m,n} x^m y^n$  is not holonomic.  $\triangleright$  Key argument: F(x, 0) has infinitely many singularities.

#### A historical example

$$a_{m,n} = \begin{cases} a_{m+1,n-2} + a_{m-2,n+1} - a_{m-1,n-1} & \text{if } m, n \ge 2\\ -\delta_{(m,n),(1,1)} & \text{if } m \le 1 \text{ or } n \le 1 \end{cases}$$

The series  $F(x, y) = \sum_{m,n \ge 2} a_{m,n} x^{m-2} y^{n-2}$  satisfies

$$(x-y^2)(y-x^2)F(x,y) = xy - G(x) - G(y)$$
, for  $G(x) = \sum_{m \ge 2} a_{m,2}x^{m+1}$ 

$$\implies x^3 - G(x) - G(x^2) = 0 \implies G(x) = \sum_{i \ge 0} (-1)^i x^{3 \cdot 2^i}$$

 $\triangleright$  G is lacunary, thus it is not holonomic, and so is F.
Two ingredients:

1. Asymptotics of excursions:

 $[t^n] F_{\mathfrak{S}}(t;0,0) \sim K \cdot \rho^n \cdot n^{\alpha},$ 

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• c and  $\rho$  are algebraic numbers depending on  $\mathfrak{S}$ 

(implied by a recent probability result [Denisov & Wachtel 2011])

2. If  $F_{\mathfrak{S}}(t; 0, 0)$  is holonomic, then it is a G-function, and  $\alpha \in \mathbb{Q}$  (implied by deep number theory results [Chudnovsky-André-Katz])

Irrationality of  $\arccos(c)/\pi$  is proven *algorithmically* in two steps:

(S1) determine the minimal polynomial,  $\mu_c$ , of c

(S2) prove that the numerator of  $\mu_c\left(\frac{x^2+1}{2x}\right)$  contains no cyclotomic polynomial factor.



▷ The algorithm proves that F(t; 0, 0) is non holonomic (and thus so is F(t; x, y)) for the 51 non-singular walks with infinite group.

The algorithm on the example



> S:=[[-1,0],[0,1],[1,0],[1,-1],[0,-1]]: > chi:=add(x^s[1]\*y^s[2],s=S);

$$\chi := \frac{1}{x} + \frac{1}{y} + x + y + \frac{x}{y}$$

> chi\_x:=numer(diff(chi,x));chi\_y:=numer(diff(chi,y));

$$\chi_x := x^2 + x^2 y - y, \qquad \chi_y := y^2 - x - 1.$$

> G:=Groebner[Basis]([chi\_x,chi\_y, numer(t<sup>2</sup>diff(chi,x,y)<sup>2</sup>/diff(chi,x,x)/diff(chi,y,y))],lexdeg([x,y],[t])): > p:=factor(op(remove(has,G,{x,y})));

$$p := (4t^2 + 1)(8t^3 + 8t^2 + 6t + 1)(8t^3 - 8t^2 + 6t - 1).$$

The polynomial p has only two real roots,  $\pm c$ . Numerical evaluation of c identifies its minimal polynomial as  $\mu_c = 8t^3 + 8t^2 + 6t + 1$ 

mu\_c:=8\*t^3+8\*t^2+6\*t+1: R:=expand(x^3\*subs(t=(x^2+1)/x/2, mu\_c),sort);

 $R(x) = x^{6} + 2x^{5} + 6x^{4} + 5x^{3} + 6x^{2} + 2x + 1.$ 

> irreduc(R),numtheory[iscyclotomic](R,x);

#### true, false

# Summary

2D classification of F(t; 0, 0) and F(t; x, y) is fully completed
robust algorithmic methods:

- Guess'n'Prove approach based on modern CA algorithms
- Creative Telescoping for integration of rational functions
- $\odot$  Brute-force and/or use of naive algorithms = hopeless.
  - E.g. size of algebraic equations for  $G(t; x, y) \approx 30$ Gb.
- Remarkable properties discovered experimentally. E.g.: all algebraic cases have solvable Galois groups

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$$G(t; 1, 1) = -\frac{3}{6t} + \frac{\sqrt{3}}{6t} \sqrt{U(t) + \sqrt{\frac{16t(2t+3)+2}{(1-4t)^2U(t)} - U(t)^2 + 3}}$$
  
where  $U(t) = \sqrt{1 + 4t^{1/3}(4t+1)^{1/3}/(4t-1)^{4/3}}.$ 

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where  $U(t) = \sqrt{1 + 4t^{1/3}(4t+1)^{1/3}/(4t-1)^{4/3}}.$ 

- lack of "purely human" proofs for many results. E.g.: non-holonomy of F(t; 0, 0) and 2F<sub>1</sub>s expressions for F(t; x, y)
- $\ensuremath{\textcircled{}}$  still missing a unified proof of: finite group  $\leftrightarrow$  holonomic
- open: is F(t; 1, 1) non-holonomic in the 51 non-singular cases with infinite group?

# Extensions

- 1. Longer 2D steps [B., Bousquet-Mélou & Melczer, in progress]
  - 680 step sets with one large step, 643 proven non holonomic, 32 of 37 have differential equations guessed.
  - 5910 step sets with two large steps, 5754 proven non holonomic, 69 of 156 have differential equations guessed.
- 2. 3D walks [B., Bousquet-Mélou, Kauers & Melczer, in progress]
  - 83 682 with 5 steps or less: B. and Kauers (2009) conjectured (up to equivalence) 35 holonomic steps. Now proved.
  - With 6 steps, 96 new holonomic cases: guessed, then proved.
  - New phenomenon (empirically discovered, no proof yet): ∃ step sets (3D Kreweras) with finite group and non-holonomic GF?!

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# Thanks for your attention!